Concurrent Systems

Nebenläufige Systeme

II. Concurrency

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October 21, 2015

Agenda

Preface

Causality

- Interdependencies
- Dimensions

Resource Sharing

- Principles
- Competition
- Synchronisation

Summary

Outline

Subject Matter

- discussion on two fundamental **abstract concepts**:

  concurrency (Ger. *Nebenläufigkeit*)
  - designates the relation of causal independent events
  - is related to events that have no mutual influence

  causality (Ger. *Kausalität*, *Ursächlichkeit*)
  - designates the relation between cause and effect
  - is the causal chain or connection of two events

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Events occur or are concurrent if none is the cause of the other.

- explanation of the relation of these concepts to **resource sharing**
  - differentiated with respect to various types of resources and sharing
  - classified as to appropriate or necessary synchronisation paradigms

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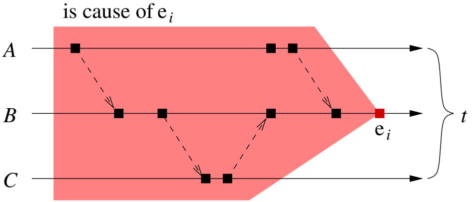
Principle of Causality

- causal chain of events related to some other event e_i:



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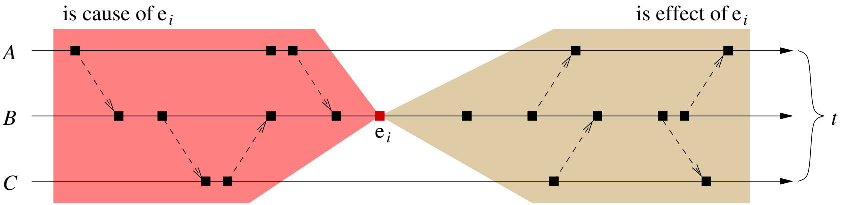

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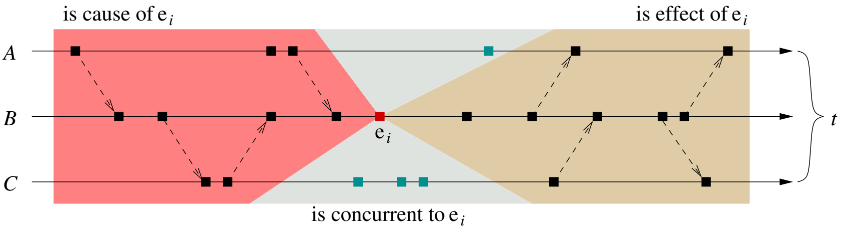

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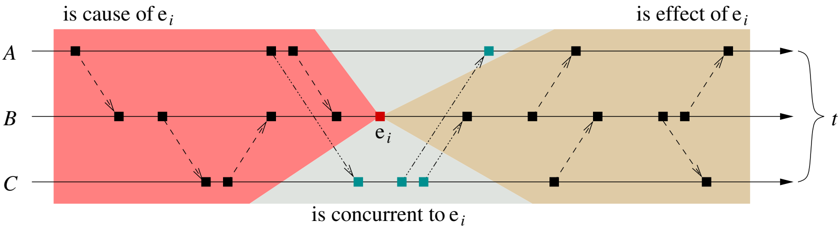

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- the event is neither cause nor effect of the other event (e_i)
  - as the case may be, it is cause/effect of other events (different from e_i) that are lying in the elsewhere (cf. dash-and-dot line)

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- interrelation of computations/events constrains concurrency

Event correlations v. Processing modes

$$\left.\begin{array}{l}\text{“is cause of”}\\ \text{“is effect of”}\end{array}\right\} \mapsto \text{sequential (realised before/at run-time)}$$

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$\hookrightarrow$ decrease of the portion of **sequential code** is an important aspect

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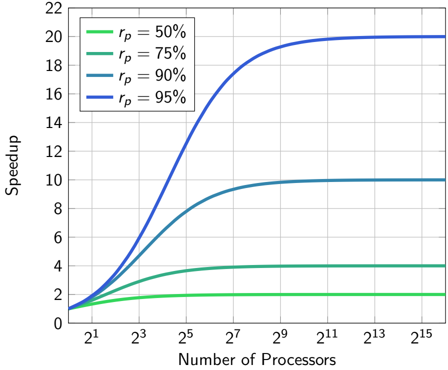
r_s ratio of sequential code

r_p ratio of parallel code, independent of n

n number of processors



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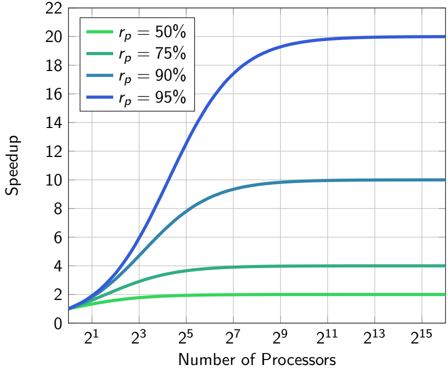

- speed-up will be constrained by **data management housekeeping**
  - the nature of this overhead appears to be sequential

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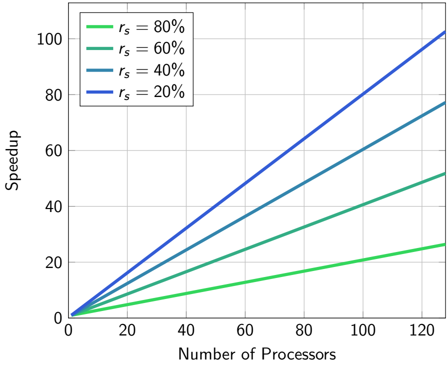

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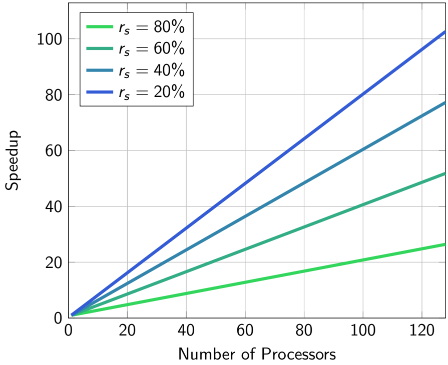

- data management housekeeping (serial part) becomes less important
  - in practise, the problem size scales with the number of processors: **HPC**[1]

[1] High Performance Computing

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- defined by the **causal order** (Ger. *Kausalordnung*) of the statements
  - as far as the logical dimension of a program is concerned
  - but there is also a **physical dimension**, namely when it comes to the execution of that program by a real processor $\leadsto$ *level of abstraction*

Level of Abstraction

- a concurrent operation (in logical terms) at a higher level can be sequential (in real terms) at a lower level

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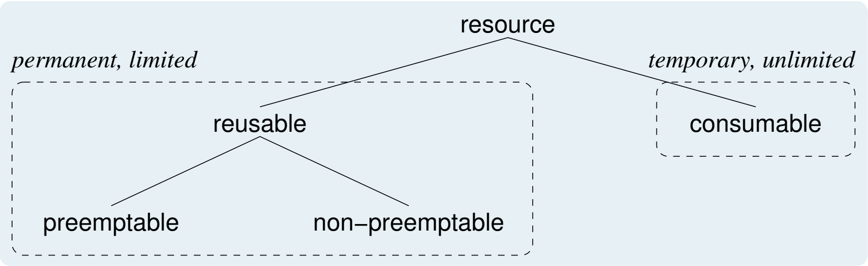
$\hookrightarrow$ reveals a *race condition*, substantial critical situation: **error**

[2] real-time processing, especially in case of hard deadlines.

Outline

Resource Classification



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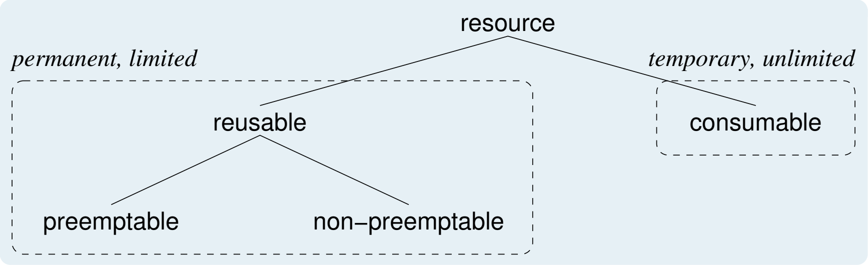

- permanent[3] resources are **reusable**, but always only of limited supply
  - they are acquired, occupied, used, and released (when no longer required)

    in-use resources: are preemptable or non-preemptable, depending on whether allocation to another occupant is possible
  - when non-preemptable, they are exclusively owned by an occupant

[3] Also referred to as "persistent".

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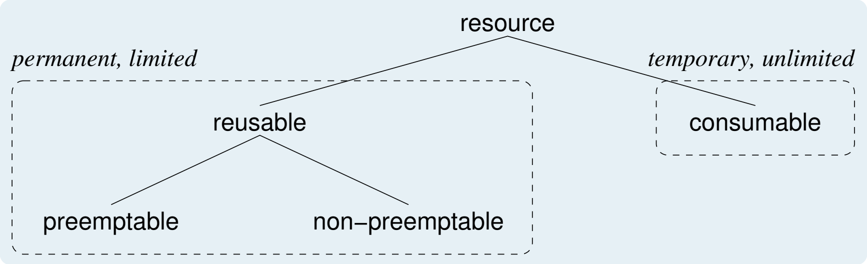

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- temporary resources are of unlimited supply, they are **consumable**
  - i.e. produced, received, used, and destroyed (when no longer required)

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Resource Peculiarities

- **hardware resources** as to be managed, e.g., by an operating system

reusable	
processor	CPU, FPU, GPU; MMU
memory	RAM, scratch pad, flash
peripheral	input, output, storage

consumable	
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- reusable and consumable resources imply different **use patterns**

Resource Use Patterns

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  - provided that the following two basic conditions (i.e., constraints) apply:
    - i resource accesses by computations may happen (quasi-) simultaneously
    - ii simultaneous accesses may cause a conflicting state change of the resource
  - simultaneous use of a **shared resource** this way must be coordinated
    - coordination may affect computations in a blocking or non-blocking manner[4]

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  - use of a **temporary resource** follows a causal course of events or actions
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- simultaneous computations **overlap** in time, interfere with each other
  - they become critical in any case if they also overlap in (identical) place

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Consolidating Example

- assuming that the following subroutines (put and get) are executed in any order and that they may also run simultaneously:

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- `put` and `get` must be subject to uni- <u>and</u> multilateral synchronisation
  - they are <u>not</u> concurrent under the assumption that was made above

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- simultaneous computations or operations, resp., are in competition:
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movl _in, %ecx
leal 1(%ecx), %eax
movl %eax, _in
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- the critical case may result in wrong reading (Ger. *Zählerwert*) of `in`/`out`
  - `in++` or `out++` are not concurrent to oneself, resp.: they are not **re-entrant**

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cf. p. 28

- assure a conflict-prone complex operation of (logical) **indivisibility**

Establishing of Synchronism

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  - ensure synchronism at **operation start**
  - lock potential overlapping out in the first place
  - synchronised operation is made of sequential code

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  - ensure synchronism at **operation end**
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- both approaches come in a variety of solutions to the same problem

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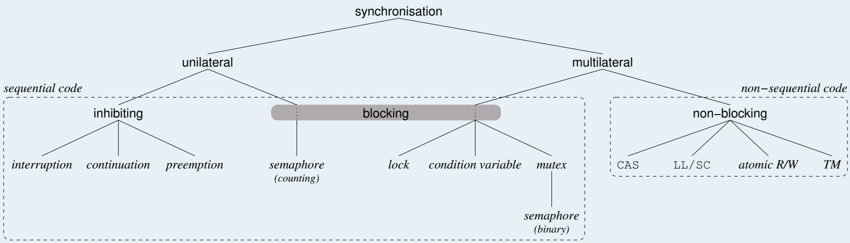
Varieties of Synchronisation



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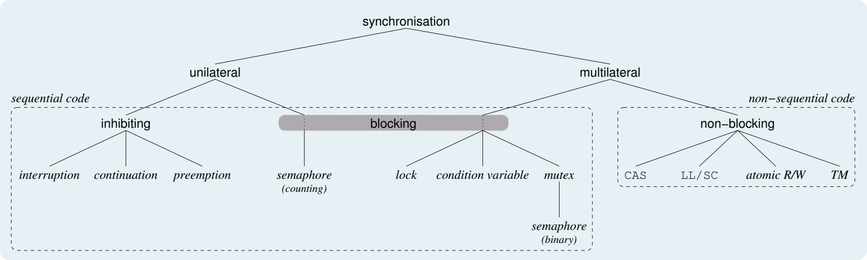

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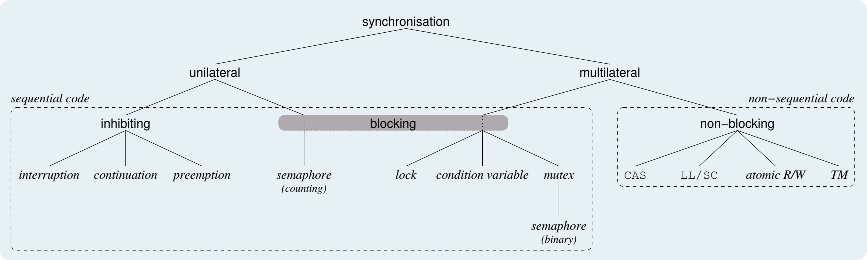

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  - bracket sequential code by a **locking protocol** ☹
  - for the most part, the original code can be reused ☺

  ↪ *pessimistic*, overlapping is <u>not</u> a rare event

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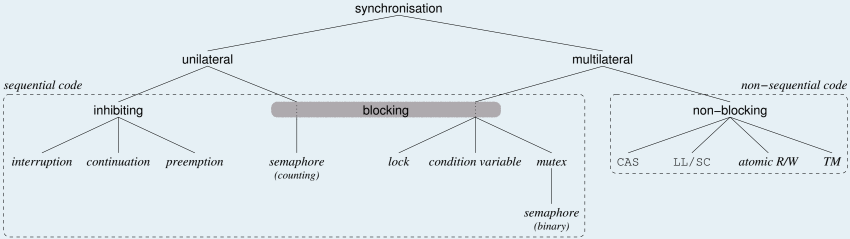

- the methods are more or less disruptive of the problematic operation:

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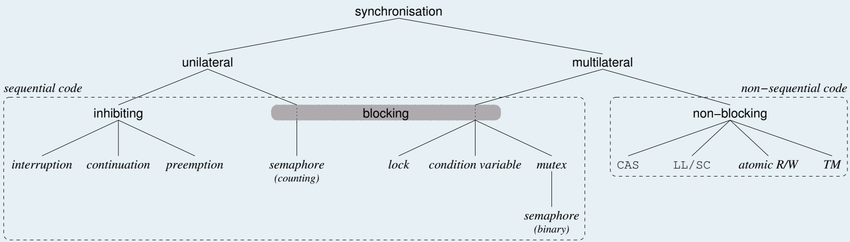

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non-sequential	■	reprogram sequential code as a **transaction**	☺
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- wherever applicable, ***downsizing sequential code*** is basic
  - i.a. Amdahl's Law (cf. p. 8) argues for non-blocking synchronization

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- delays computations subject to resource availability
  - takes effect only in case of co-occurrence (overlapping)
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- it bears repeating: ***downsizing sequential code*** is basic
  - where possible, non-blocking synchronisation should be the first choice
- but even then: there is no all-in-one approach for every purpose. . .

Outline

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  - but simultaneity my also cause overlapping that must be constrained
- synchronism ensures that overlapped complex operations do right
  - the individual sub-steps will be strictly executed *interim* (consecutively) *or*
  - a *transaction* will take care for consistent (pseudo-) parallel execution

Relativity of Simultaneity

- the concept of (distant) simultaneity is not absolute, but depends on the **frame of reference** (Ger. *Bezugssystem*) an observer takes
  - moving- and fixed-platform thought experiment [2, p. 768]:

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    - while multiple invocations of the former will take place sequentially,[7] the corresponding ones of the latter may come about non-sequentially
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  - operations must be resolved **cross-level** (from "fixed platform" observed) in order to realise their ability for concurrency or need for synchronism

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Résumé

- computations can be **concurrent** if none needs a result of the other
  - they must be free of data and control-flow dependencies
- in order to be concurrent, computations must be **simultaneous**
  - quasi-simultaneous through partial virtualization (hardware multiplexing) or real simultaneous by multiprocessing (hardware multiplication)
  - both techniques will induce computations to overlap in time and place
- **overlapping** in time cause interference but is the lesser of two evils
  - more critical is overlapping in place relating to the same resource
  - particularly with regard to the same (i.e., shared) memory area
- critical overlapping must be counteracted through **synchronisation**
  - i.e., coordination of the cooperation and competition between processes
  - here: uni- or multilateral synchronisation, depending on the resource type
- synchronisation ensures for **indivisibility** of a computation cycle
  - at the outset: physical, in blocking manner, by being pessimistic ☹
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Consolidating Example Revisited

cf. p. 16

- **bounded buffer** using a counting semaphore [3] for unilateral and an ELOP (x86) for multilateral synchronisation

```c
typedef int semaphore_t;

extern void P(semaphore_t*);
extern void V(semaphore_t*);

semaphore_t free = 80;
semaphore_t empty = 0;

static inline int fai(int *ref) {
  int aux = 1;

  asm volatile("lock; xaddl %0,%1"
    : "=r" (aux), "=m" (*ref)
    :  "0" (aux),  "m" (*ref));

  return aux;
}
char buffer[80];
unsigned in = 0, out = 0;

void put(char item) {
  P(&free);
  buffer[fai(&in) % 80] = item;
  V(&empty);
}

char get() {
  char item;

  P(&empty);
  item = buffer[fai(&out) % 80];
  V(&free);

  return item;
}
```

free
- controls the number of unused buffer entries
- *P* prevents from buffer overflow, *V* signals reusable resource

empty
- controls the number of used buffer entries
- *P* prevents from buffer underflow, *V* signals consumable resource

fai
- indivisibly *fetch and increment* specified counter variable