

Concurrent Systems

Nebenläufige Systeme

XII. Transactional Memory

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Agenda

Preface

Principles

- General

- Characteristic

- Operation

Utilisation

- Abstraction

- Exemplification

- Discussion

Summary



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 - explicitly versus implicitly transactional approaches
 - hardware (HTM), software (STM), or hybrid (HyTM) solutions



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 - load, store, and commit for the explicit case
 - begin and end for the implicit case
 - abort for both cases—the exception proves the rule. . .



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 - failure of transactions: frequency, reason, alternative measures
 - type of synchronisation: unilateral, multilateral



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Universal Remedy?

The bigger the critical shared state is, the better TM seems to be. But what about support, overhead, control, and coordination?





Source: http://de.wikipedia.org/wiki/Transzendentele_Meditation





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 - similar to blocking synchronisation such as mutual exclusion, the secured program sections describes a seemingly sequential process
 - but unlike that synchronisation pattern, the respective program sections may be run by non-sequential processes
 - within those sections, on a simple load/store basis, instructions are given what data need to be kept consistent
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- in spite of everything, not run the **risk** of overstating sequential and understating parallel thinking. . .



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Warm-Up

- to come to the point, TM has a silver lining but also a demerit



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 - pros
 - allows for the definition of customized atomic operations that apply to a group of possibly arbitrary computer words
 - can be seen as a caching method for implementing data structures in a lock-free manner [2]
 - technically feasible as “straightforward extensions to multiprocessor cache-coherence protocols” [10]
 - offers a more convenient handling compared to, e.g., a multi-word CAS using (software-implemented CAS-based) LL/SC [14]



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- cons**
- may be a replacement for multilateral synchronisation (i.e., mutual exclusion using e.g. locks or binary semaphores), only
 - neither facilitates nor supports, but rather hampers, unilateral (i.e., logical/conditional) synchronisation
 - prone to overhead in case of mindless reuse of external functions or procedures from libraries, for instance
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 - tempt developers of non-sequential programs to see things through rose-coloured glasses
- TM is **a means to an end**—and is far from being a cure-all. . .



Levels of Abstraction

- depending on the rootedness of the implementation, divided into:
 - HTM ■ **hardware** transactional memory [10]

STM ■ **software** transactional memory [17]

HyTM ■ **hybrid** transactional memory [5]



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 - memory instructions indicate each single transactional load/store
 - may also provide instructions to start and commit transactions
 - *implicitly transactional* (SLE [16], Rock [18], PPC [7], HLE [12])
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- STM
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 - heterogeneous transactions: (1) HTM-based and (2) STM-based
 - (1) gets aborted if in conflict with (2), may be restarted as (2)
- howsoever, **linguistic support** is desirable—but, with or without it, TM is no panacea to solve all non-sequential programming issues



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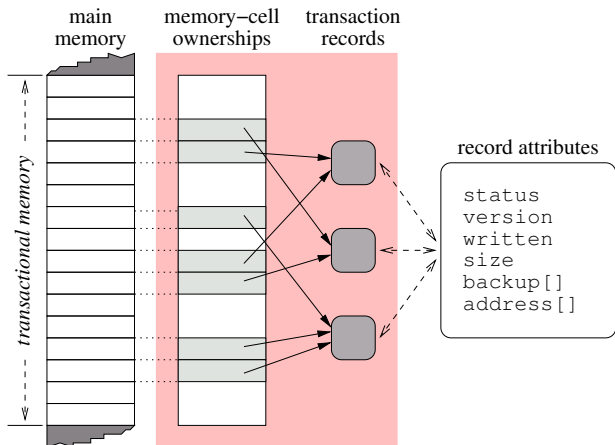


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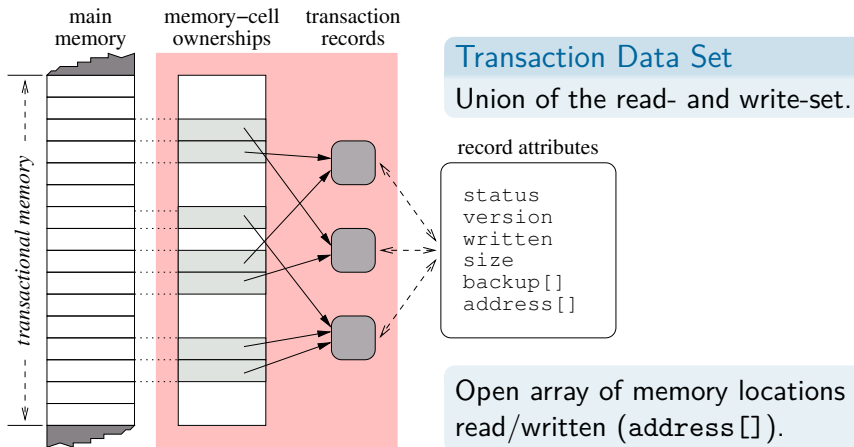
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- but, not yet really common in available processors architectures:
 - IBM Blue Gene/Q (PowerPC A2 [7])
 - limited to multi-versioned L2 cache (20 MiB out of 32 MiB, [19, p. 129])
 - “watch granule” is 64 B [11, p. 509], same as cache line size
 - Intel Haswell (TSX [12])¹
 - transaction size limited to L1 cache (64 KiB), 64 B cache line

¹Mindless of the TSX bug [13, p. 47], which leaves TSX barred for normal use.



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- a **generation** number (preferable 64-bit) makes a transaction unique
 - additional parameter of steps 1, 2.1, 2.3, and 2.4: needs to be checked



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- an **implicitly transactional** model is assumed, but not the only way
 - by concept, an explicitly transactional approach is feasible as well



Minimal Subset of System Functions

- operations for accessing memory (implicitly or explicitly):
 - load* ■ transfers a value from shared memory to a private placeholder
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 - ii read any location in the current write set



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- further operations are customary, according to circumstances
 - depending on the level of abstraction the TM system is associated with



- it is assumed that contention of simultaneous processes is improbable
 - a **successful commit** seems to be probable, other than abort and retry

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1 extern word_t foo, bar, foobar;  
2 do {  
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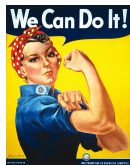


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 - context switches (e.g., system calls, process dispatching)
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- thus, be aware of reasons and frequency of the failure of transactions
 - if applicable, take care of region-specific **counteractive measures**
 - reflect on alternative concepts/solutions in achieving data consistency



- a more advanced abstraction is to merely declare an **atomic region**
 - at the expense of a loss of control of the extent of the actual data set

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1 extern word_t foo, bar, foobar;  
2 begin(&&dropout);  
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- **retry-loop concealment** is not always an advantageous measure
 - aside from other exceptional events (p. 15), retries are due to contention
 - contention control depends not only on dynamic but also static data
 - i.e., number of contending processes and duration of a single retry
 - whereby the latter is determined by the regions's execution path length
 - *begin/end* are unaware of **expectable execution times** of atomic regions



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```
1 extern void btx(void *);           /* begin */
2 extern void atx();                 /* abort */
3 extern bool ctx();                 /* commit */
4 extern void etx();                 /* end */
5
6 extern long ltx(void *);           /* load */
7 extern void stx(void *, long);     /* store */
```

- in case of STM, it is worth to consider the following **refinements**:
 - upper-bound size of the read- and write-set in btx
 - specification of the reason of abort in abx
 - declaration of further modes of operation (flags) in btx
 - additional (first) parameter indicating this transaction in each operation
- however, as (most of) these depend on the program structure of the **transactional region**, determination should be up to the compiler



LIFO-List Revisited I

```
1 inline void push_dos(stack_t *this, chain_t *item) {  
2     item->link = this->head.link;  
3     this->head.link = item;  
4 }
```



LIFO-List Revisited I

```
1 inline void push_dos(stack_t *this, chain_t *item) {
2     item->link = this->head.link;
3     this->head.link = item;
4 }

5 void push_tm_it(stack_t *this, chain_t *item) {
6     btx(0);
7     item->link = this->head.link;
8     this->head.link = item;
9     etx();
10 }
```



LIFO-List Revisited I

```
1 inline void push_dos(stack_t *this, chain_t *item) {
2     item->link = this->head.link;
3     this->head.link = item;
4 }
```

```
5 void push_tm_it(stack_t *this, chain_t *item) {
6     btx(0);
7     item->link = this->head.link;
8     this->head.link = item;
9     etx();
10 }
```

`item->link`

Unnecessary
data-set member.



LIFO-List Revisited I

```
1 inline void push_dos(stack_t *this, chain_t *item) {
2     item->link = this->head.link;
3     this->head.link = item;
4 }

5 void push_tm_it(stack_t *this, chain_t *item) {
6     btx(0);
7     item->link = this->head.link;
8     this->head.link = item;
9     etx();
10 }

11 void push_tm_et(stack_t *this, chain_t *item) {
12     do {
13         item->link = (chain_t *)ltx(&this->head.link);
14         stx(&this->head.link, (long)item);
15     } while (!ctx());
16 }
```

`item->link`

Unnecessary
data-set member.



FIFO-List Revisited I

```
1 inline void chart_dos(queue_t *this, chain_t *item) {
2     item->link = 0;           /* finalise chain */
3     this->tail->link = item;  /* append item */
4     this->tail = item;       /* set insertion point */
5 }
```



FIFO-List Revisited I

```
1 inline void chart_dos(queue_t *this, chain_t *item) {
2     item->link = 0;           /* finalise chain */
3     this->tail->link = item;  /* append item */
4     this->tail = item;       /* set insertion point */
5 }
6 void chart_tm_it(queue_t *this, chain_t *item) {
7     item->link = 0;
8     btx(0);
9     this->tail->link = item;
10    this->tail = item;
11    etx();
12 }
```



FIFO-List Revisited I

```
1 inline void chart_dos(queue_t *this, chain_t *item) {
2     item->link = 0;                /* finalise chain */
3     this->tail->link = item;       /* append item */
4     this->tail = item;            /* set insertion point */
5 }

6 void chart_tm_it(queue_t *this, chain_t *item) {
7     item->link = 0;
8     btx(0);
9     this->tail->link = item;
10    this->tail = item;
11    etx();
12 }

13 void chart_tm_et(queue_t *this, chain_t *item) {
14     item->link = 0;
15     do {
16         stx(&this->tail->link, (long)item);
17         stx(&this->tail, (long)item);
18     } while (!ctx());
19 }
```

Both TM variants appear to be equivalent.



LIFO-List Revisited II

```
1 chain_t *wear_dos(stack_t *this) {
2     chain_t *node = this->head.link;
3     this->head.link = 0;
4     return node;
5 }
```



```
1 chain_t *wear_dos(stack_t *this) {
2     chain_t *node = this->head.link;
3     this->head.link = 0;
4     return node;
5 }

6 chain_t *wear_tm(stack_t *this) {
7     chain_t *node;
8     do {
9         node = ltx(&this->head.link);
10        stx(&this->head.link, 0);
11    } while (!ctx());
12    return node;
13 }
```



```
1 chain_t *wear_dos(stack_t *this) {
2     chain_t *node = this->head.link;
3     this->head.link = 0;
4     return node;
5 }

6 chain_t *wear_tm(stack_t *this) {
7     chain_t *node;
8     do {
9         node = ltx(&this->head.link);
10        stx(&this->head.link, 0);
11    } while (!ctx());
12    return node;
13 }

14 chain_t *wear_wfs(stack_t *this) {
15     return FAS(&this->head.link, 0);
16 }
```

Overshoot

Definitely, TM is no magic bullet...



All that Glitters is not Gold...

The TM programming model itself, whether implemented in hardware or software, introduces complexities that limit the expected productivity gains, thus reducing the current incentive for migration to transactional programming and the justification at present for anything more than a small amount of hardware support. [4, p. 55]



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- **logical/conditional synchronisation**, e.g. condition variables [6]:
 - waiting on a condition inside a transaction is difficult or impossible
 - difficult, e.g., in case of an I/O operation that cannot be rolled back
 - impossible, if the transactional process is implemented as kernel-level thread²

²Assuming that TM applies to user-level processes, only—which is usual.



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 - impossible, if the transactional process is implemented as kernel-level thread²
 - as a signalling transaction may abort, the stated condition never occurred
 - furthermore, signaller and signallee transactions may happen simultaneously, which is prone to lost-wakeup as the latter may complete before the former

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 - as a signalling transaction may abort, the stated condition never occurred
 - furthermore, signaller and signallee transactions may happen simultaneously, which is prone to lost-wakeup as the latter may complete before the former
- thus, TM is merely an abstraction to **multilateral synchronisation**
 - most attractive semantics is its “single global lock atomicity” [3] ☹️

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Outline

Preface

Principles

- General

- Characteristic

- Operation

Utilisation

- Abstraction

- Exemplification

- Discussion

Summary



- TM abstractions as to the different rootedness of the implementation
 - HTM ■ hardware, explicitly/implicitly transactional, hardly available
 - STM ■ software, lock-less/based solutions, metadata overhead
 - HyTM ■ hybrid, try HTM first, fall back on STM in critical situations
- principle concepts of TM and functions or instructions, respectively
 - read set, write set, and the union thereof: data set
 - load, store, commit, abort, begin, end—and more. . .
- examination and discussion of the pros and cons of TM
 - especially limited hardware support still hampers wide use
 - independent thereof, programming introduces other types of complexities
 - also because it merely is an abstraction to multilateral synchronisation
- TM is a means to an end, it has a silver lining but also a demerit. . .

Transactional Memory Should Be an Implementation Technique, Not a Programming Interface. [3]



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```
1 typedef struct ref {
2     int *label; /* actual location in (shared) memory */
3     int owner; /* reservation number: initial anything but -1 */
4 } ref_t;
5
6 inline int ll(ref_t *ref, int key) {
7     int owner, value;
8     do {
9         owner = ref->owner;
10        value = *(ref->label);
11    } while ((ref->owner == -1) || !CAS(&ref->owner, owner, key));
12    return value;
13 }
14
15 inline bool sc(ref_t *ref, int key, int val) {
16     bool done;
17     if ((done = CAS(&ref->owner, key, -1))) {
18         *(ref->label) = val;
19         ref->owner = 0;
20     }
21     return done;
22 }
```

- key ■ unique transaction number
- advanced when transaction completes



LIFO-List Revisited III

```
1 inline chain_t *pull_dos(stack_t *this) {
2     chain_t *node;
3     if ((node = this->head.link))
4         this->head.link = node->link;
5     return node;
6 }

7 chain_t *pull_tm(stack_t *this) {
8     chain_t *node;
9     do {
10         if ((node = (chain_t *)ltx(&this->head.link)))
11             stx(&this->head.link, (long)node->link);
12     } while (!ctx());
13     return node;
14 }
```

- the implicitly transactional variant would unnecessarily include *node* in the transaction data set...



FIFO-List Revisited II

```
1 inline chain_t* fetch_dos(queue_t *this) {
2     chain_t *node;
3     if ((node = this->head.link)           /* filled? */
4         && !(this->head.link = node->link)) /* last item? */
5         this->tail = &this->head;        /* reset */
6     return node;
7 }
8 chain_t* fetch_tm(queue_t *this) {
9     chain_t *node;
10    do {
11        if ((node = (chain_t *)ltx(&this->head.link))) {
12            stx(&this->head.link, (long)node->link);
13            if (!node->link)
14                stx(&this->tail, (long)&this->head);
15        }
16    } while (!ctx());
17    return node;
18 }
```

■ the implicitly transactional variant would unnecessarily include *node* in the transaction data set...



FIFO-List Revisited III

```
1 inline chain_t *drain_dos(queue_t *this) {
2     chain_t *head = this->head.link;
3     this->head.link = 0;          /* null item */
4     this->tail = &this->head;    /* linkage item */
5     return head;
6 }

7 chain_t *drain_tm(queue_t *this) {
8     chain_t *head;
9     do {
10         head = (chain_t *)ltx(&this->head.link);
11         stx(&this->head.link, 0);
12         stx(&this->tail, (long)&this->head);
13     } while (!ctx());
14     return head;
15 }
```

- the implicitly transactional variant would unnecessarily include *head* in the transaction data set...

