

## E Distributed Objects in CORBA

- ★ Motivation
- Location
  - ◆ Transparency of location
- Heterogeneity
  - ◆ Different hardware,
  - ◆ different operating systems, and
  - ◆ different programming languages are used in distributed systems.
  - ◆ Transparency of heterogeneity
- Services
  - ◆ Name server
  - ◆ Time server
  - ◆ ...

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## E.2 CORBA Overview

### 1 Design Goals

- CORBA is a standard
  - ◆ There are standard documents
  - ◆ Definition of CORBA compliance
  - ◆ Vendors build implementations of the standard (e.g., VisiBroker, OmniBroker, Orbix, etc.)
- CORBA wants to abstract from
  - ◆ hardware,
  - ◆ operating system, and
  - ◆ programming language.

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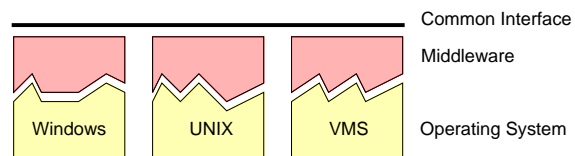
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## E.1 Middleware Approach

- Middleware: a piece of software between operating system and application



- ◆ Middleware provides services for distributed programming
- **We need:** middleware for distributed object-based programming
- ★ CORBA – Common Object Request Broker Architecture, a standard of the OMG

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## 1 Design Goals (2)

[E.2 CORBA Overview](#)

- Interoperability
  - ◆ An application shall run on different CORBA implementations without major changes
  - ◆ Applications on different CORBA implementations shall be able to communicate
- Embedding of legacy applications
  - ◆ Legacy applications can be encapsulated and can act as CORBA objects
- Vision of business objects or components
  - ◆ CORBA objects represent all kind of data and proceedings of a company

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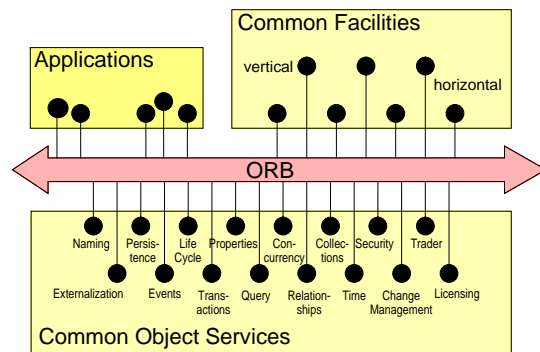
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## 2 Architectural Overview

E.2 CORBA Overview



- Object Management Architecture – OMA

## 3 CORBA Implementations

E.2 CORBA Overview

- A CORBA implementation must contain
  - ◆ the implementation of the core architecture
  - ◆ one language mapping (e.g., for C++)
- A CORBA implementation may contain
  - ◆ an arbitrary number of services
  - ◆ functionality for interoperability with other CORBA implementations (GIOP/IOP)
- The implementation is free **how** it realizes the requirements of the standard
  - ◆ Many different implementations are possible
    - Daemon-based
    - Library-based
    - ...

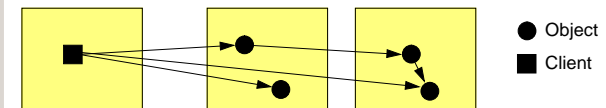
## 2 Architectural Overview (2)

E.2 CORBA Overview

- Application objects
  - ◆ Whole object approach for distributed objects
  - ◆ Client/server architecture
- Object Request Broker (ORB)
  - ◆ Communication backbone
  - ◆ Enables objects to communicate
- CORBA Services
  - ◆ Basic services for distributed programming
  - ◆ Extends the ORB's functionality
  - ◆ Services look like objects
- CORBA Facilities
  - ◆ Application-specific services

## E.2 Application Objects

- Distributed Objects
  - ◆ Identity
  - ◆ State
  - ◆ Operations (methods)
  - ◆ CORBA objects can be invoked (implement server side)
  - ◆ CORBA objects may act as clients



- Clients need not to be CORBA objects
  - ◆ CORBA clients can be processes, etc.

## E.2 Application Objects (2)

- Distributed objects form an application
  - ◆ Objects cooperate by communicating with each other
- ★ Example: printer management system
  - ◆ Client objects
  - ◆ Spooler objects
  - ◆ Printer objects
- "Whole object approach"
  - ◆ Client-side stub can contact the server object

## 1 Interface Definition Language (2)

### ■ Example

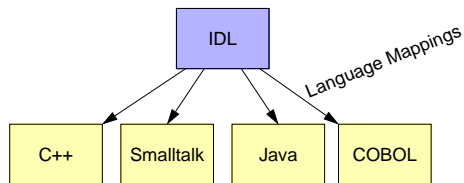
```
module MyModule
{
  interface MyInterface
  {
    attribute int lines;
    void printLine( in string toPrint );
  };
};
```

IDL to C++

```
namespace MyModule {
  class MyInterface : ... {
  public:
    virtual CORBA::Long lines();
    virtual void lines( CORBA::Long _val );
    void printLine( const char *toPrint );
    ...
  };
};
```

## 1 Interface Definition Language (IDL)

- Language for describing object interfaces
  - ◆ Independent of implementation language of an object
  - ◆ Language mapping defines how IDL primitives are mapped to the concepts of a specific programming language
  - ◆ Language mapping is part of the CORBA standard
  - ◆ Language mappings exist for C, C++, Smalltalk, COBOL, and Java
  - ◆ IDL is similar to C++



## 1 Interface Definition Language (3)

### ■ Example

```
module MyModule
{
  interface MyInterface
  {
    attribute int lines;
    void printLine( in string toPrint );
  };
};
```

IDL to Java

```
public interface MyModule.MyInterface extends ... {
  public int lines();
  public void lines( int lines );
  public void printLine( java.lang.string toPrint );
  ...
};
```

## 1 Interface Definition Language (4)

### ★ Advantages

- ◆ Transparency of implementation language
- ◆ Enables interoperability

### ▲ Disadvantages

- ◆ Object interface must be defined in the target language **and** in IDL
- ◆ IDL is quite expressive
  - Language that do not provide the right mechanisms will have a complex language mapping
- ◆ If a language has special features they cannot be used because they are not part of IDL

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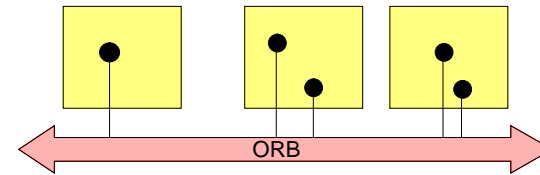
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## E.3 Object Request Broker – ORB



- The ORB is the communication backbone of a CORBA implementation
- All communication is handled by the ORB
  - ◆ ... within an address space / process
  - ◆ ... between address spaces / processes
  - ◆ ... between address spaces / processes of different ORBs
- The ORB implements location transparency

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## 2 Process of Creation and Binding

### ■ Creation of a server object

- ◆ Description of the object interface in IDL
- ◆ Programming the server object in an implementation language
- ◆ Registration of the object at the ORB (the OA respectively)
  - The ORB creates an Interoperable Object Reference (IOR)

### ■ Binding to the server object at client side

- ◆ Retrieval of an object reference to the server
  - Result of a name-server query
  - Return parameter of a method call
  - Retrieval from outside of the system: user knows the reference as a string (IORs can be converted to strings and back)
- ◆ Creation of the client stub
- ◆ Method invocations using the stub object

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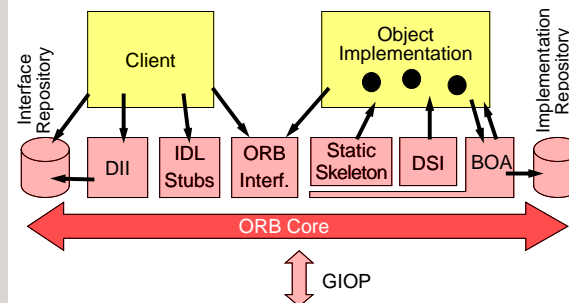
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## E.3 Object Request Broker (2)

### ■ Central components of an ORB



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## 1 Static Stubs

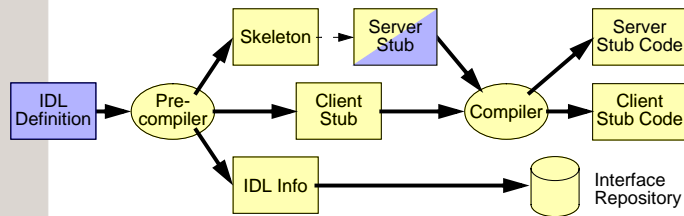
- Stubs on client and server side
  - ◆ As soon as we exactly know the interface of an object, we can create static stubs from it.
  - ◆ Static stubs are automatically created from the IDL description
  - ◆ Server side stubs are called *skeletons*
    - CORBA skeletons have to be filled with the actual implementation
- Tasks of the stubs
  - ◆ Marshalling of parameters
  - ◆ Send and receive of request and reply messages using the ORB core

## 2 Interface Repository

- Data base for interface definitions in IDL
- Query the data base for
  - ◆ Type checking
    - Does the stub's type match the object's type?
    - Inter-ORB operations: insertion into multiple repositories
  - ◆ Retrieval of meta data by clients and tools
    - Dynamic invocations
    - Debugging
    - Class browser
  - ◆ Implementing method *get\_interface* of each object
- Writing to the data base using
  - ◆ IDL compiler
  - ◆ Write methods

## 1 Static Stubs (2)

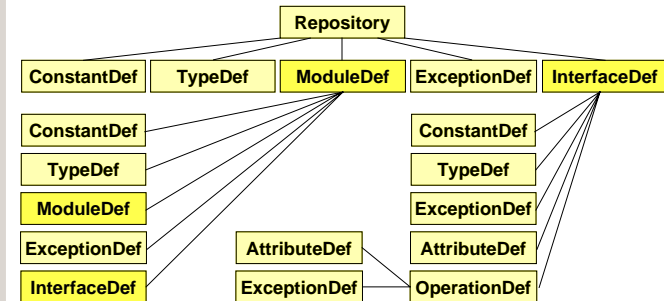
- Stub creation



- ◆ From IDL, a precompiler generates a number of output files.
- ◆ The skeleton contains a frame for filling in the actual methods of the object implementation. The skeleton already contains the server-stub functionality.
- ◆ The client-stub file contains the code for the client stub.
- ◆ An additional file may contain data for the interface repository.
- ◆ Usually the stubs have to be compiled by the target-language compiler.

## 2 Interface Repository (2)

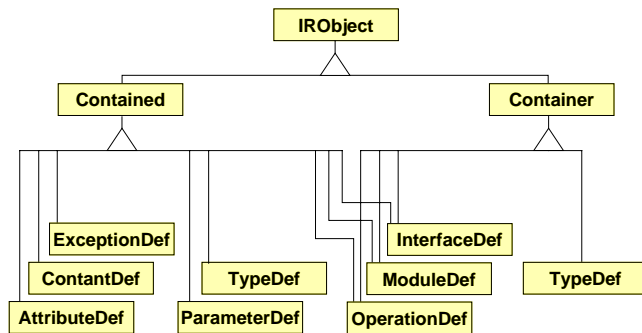
- Stored information / composition of an IDL file



- ◆ Hierarchy of the components of an IDL file, and of the interface repository respectively

## 2 Interface Repository (3)

- Inheritance hierarchy of the IDL interfaces of IR objects

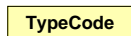


## 3 Dynamic Invocation Interface (DII)

- DII allows the invocation of methods whose object interface was not known at compile time
  - Interface description can be retrieved using the interface repository
- Single steps of an invocation (which have to be mapped by the language binding)
  - Retrieve signature of the method from the interface repository
  - Create list of parameters
  - Compose request
  - Invoke method
    - with RPC (synchronous call)
    - with asynchronous RPC (asynchronous call)
    - with datagram message (no reply, no reliability)

## 2 Interface Repository (4)

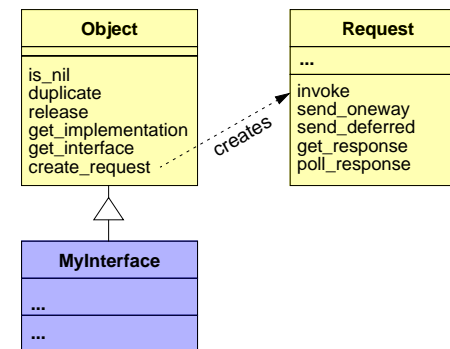
- Standard types are represented by so-called *type codes*



- Representation of basic types: **int, float, boolean**
- Representation of composite standard types: **union, struct, enum**
- Representation of template types and complex declarators: **sequence, string, array**
- Type codes represent type and structure as objects with an standard IDL interface
  - Operation for comparing type codes
  - Operation for obtaining the description of the type

## 3 Dynamic Invocation Interface (2)

- IDL definitions for the DII



## 4 Dynamic Skeleton Interface (DSI)

- DSI allows a CORBA server to accept method invocations for objects whose interfaces are only dynamically retrievable, e.g., in
  - ◆ Bridges to other ORBs
  - ◆ CORBA-encapsulated data bases
  - ◆ Dynamically created objects and interfaces
- The DSI identifies the object for which the invocation is made for
  - ◆ A call-back function is invoked that gets all the necessary information
  - ◆ The call-back function has to invoke the corresponding object

### ★ Note:

DII and DSI are not distinguishable from static stubs, i.e., they are fully interoperable with static stubs.

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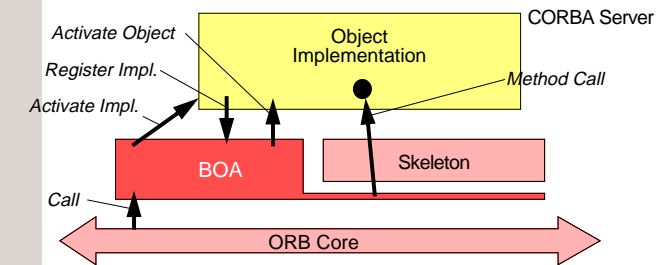
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## 6 Basic Object Adaptor (BOA)

- Activation and deactivation of CORBA servers
  - ◆ Implementation of an objects is not always active in the sense of the operating system



- Note: Do not mix up server objects with CORBA servers.

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## 5 Object Adaptor

- The object adaptor is a local representative of ORB services
  - ◆ It generates object references (for new objects)
  - ◆ It maps object references to implementations
  - ◆ It forwards incoming method calls
  - ◆ It authenticates the caller (implementing CORBA's security functionality)
  - ◆ It activates and deactivates the object and its implementation
  - ◆ It registers server classes within the implementation repository
- CORBA defines the mandatory Basic Object Adaptor
  - ◆ Basic functionality
- Another example: an OODB Adaptor
  - ◆ Connecting an object-oriented data base to CORBA
  - ◆ The OODB Adaptor converts all objects in the data base to CORBA objects

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## 6 Basic Object Adaptor (2)

- Activation strategies
  - ◆ **Shared Server**  
A server process contains multiple objects of an implementation
  - ◆ **Persistent Server**  
Activation strategy is realized outside of the BOA; server is always running (persistent); server process hosts multiple objects of an implementation
  - ◆ **Unshared Server**  
One server process per object instance
  - ◆ **Server per Method**  
One server process per method invocation

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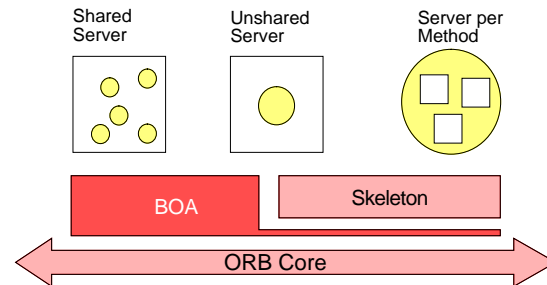
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## 6 Basic Object Adaptor (3)

### ■ Activation strategies



- ◆ Circles are objects
- ◆ Squares are CORBA servers

## 8 Inter-ORB Communication

### ■ GIOP – General Inter-ORB Protocol

- ◆ Basic protocol for the interaction of two ORBs
- ◆ Common Data Representation (CDR) converts IDL-compliant parameters into a serial byte stream
- ◆ IIOP – Internet Inter-ORB Protocol (GIOP over TCP/IP; implementation is mandatory)
- ◆ Other implementations of GIOP are possible

### ■ ESIIOP – Environment Specific Inter-ORB Protocols, e.g., DCE/ESIIOP

- ◆ Inter-ORB protocol on the basis of DCE RPC

### ■ IOR – Interoperable Object Reference

- ◆ String representation
- ◆ Several profiles (one of it is IIOP)

## 7 Implementation Repository

### ■ Data base for object implementations

- ◆ Data to localize and activate object implementations
- ◆ Information for debugging and administration
- ◆ ...

### ■ Implementation repository is highly implementation dependent.

### ■ Query the data base

- ◆ Implementation of the method *get\_implementation* of each object

### ■ Writing the data base by

- ◆ External tools
- ◆ The BOA at the start of a server

## E.4 Services

### ■ Basic services of a distributed system as extensions of the ORB

- ◆ Naming service
- ◆ Transaction service
- ◆ Persistent object service
- ◆ Life cycle service
- ◆ ...

### ■ Services are accessible by IDI interfaces

- ◆ No additional primitives necessary; method invocation is enough

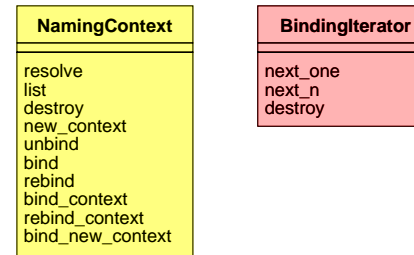
# 1 Naming Service

- CORBA defines an hierarchical naming service similar to the UNIX file system
  - ◆ Name space forms a tree
  - ◆ Names consist of multiple components (syllables)
    - e.g., < "usr"; "home"; "myobject" >
    - (UNIX: "/usr/home/myobject" )
  - ◆ Interface to the naming service is defined by IDL

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# 1 Naming Service (3)

- Interface of a Naming Context and the Binding Iterator

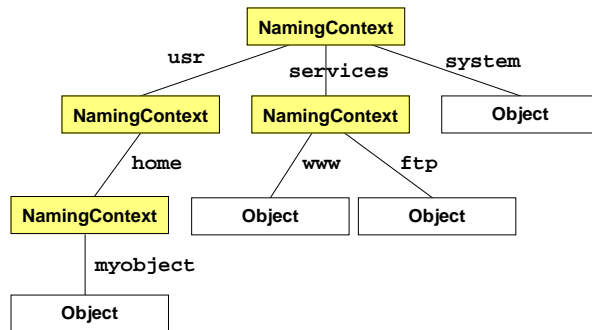


- ◆ Naming Context acts like a directory
- ◆ Objects or other contexts can be bound to a name
- ◆ Listing the context returns an array and an iterator that allows you to look at the next entries

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# 1 Naming Service (2)

- Example of a name tree



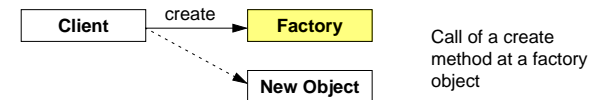
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# 2 Life Cycle Service

- Controls the life cycle of an object
  - ◆ Creation
  - ◆ Copying, Migration
  - ◆ Deletion
- Life Cycle Service defines the common interface for controllable objects

## ★ Modell des Lebenszyklus

- ◆ Creation of an Object:

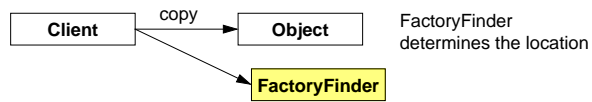


- Protocol and interface of the factory is not defined and probably dependent on the kind of object to create

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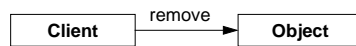
## 2 Life Cycle Service (2)

### ◆ Copying and migration:



- Locations can be a specific computer or a group of computers, etc.

### ◆ Deletion:



- With the call of remove the object is deleted.

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## 2 Life Cycle Service (4)

### ■ Example: implementation of the copy meethod

- ◆ Client invokes copy method and passes a reference to a FactoryFinder
- ◆ The copy method is provided by the programmer **or** by the CORBA implementation
- ◆ The copy methode calls the FactoryFinder and selects a suitable factory that is used to create a new object
- ◆ The new object will be initialized with the data of the existing object

### ▲ To the outside: clearly defined interface

### ▲ To the inside: open and implementation dependent, e.g., protocol between copy method and factory

### ■ In CORBA 2.0 relationship between objects is considered for copying and mobility: part-of relationship, etc.

- ◆ Parts of the object are also copied or migrated (deep copy vs. shallow copy)

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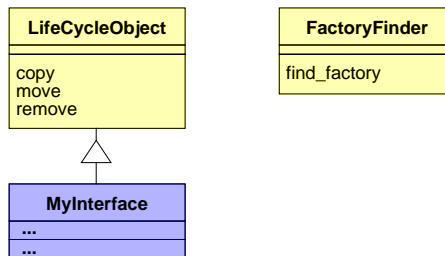
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## 2 Life Cycle Service (3)

### ■ Objects need to implement the LifeCycleObject interface



- ◆ copy and move need a FactoryFinder object to find a factory object which in turn will create the copy or migrated object

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